

# Measuring the Impact of Adversarial Errors on Packet Scheduling Strategies

A. Fernandez Anta<sup>1</sup>, C. Georgiou<sup>2</sup>, D. R. Kowalski<sup>3</sup>, J. Widmer<sup>1</sup>, E. Zavou<sup>1,4</sup>  
<sup>1</sup>Institute IMDEA Networks, <sup>2</sup>University of Cyprus, <sup>3</sup>University of Liverpool, <sup>4</sup>Universidad Carlos III de Madrid

## Introduction

**Problem definition.** Packet transmission over one unreliable link with:

- worst occurrence of link errors
- adversarial or stochastic packet arrivals
- instantaneous or deferred feedback



**Efficiency metric. Relative throughput:**

Long-term competitive ratio [1] with respect to the throughput achieved by the optimal offline algorithm.

$L_X(A, E, t)$ : total length of packets transmitted successfully by an algorithm  $X$  up to time  $t$ , under packet arrival pattern  $A$  and error pattern  $E$ .

$T_{Alg}(A, E, t)$ : partial relative throughput of algorithm  $Alg$  up to time  $t$  and under  $A$  and  $E$ .

$$T_{Alg}(A, E, t) = \frac{L_{Alg}(A, E, t)}{L_{OPT}(A, E, t)}$$

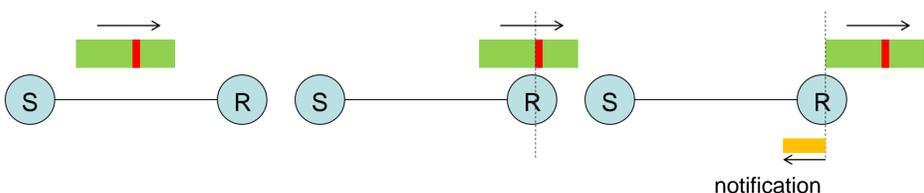
$T_{Alg}$ : relative throughput of  $Alg$  in the long-term ( $t \rightarrow \infty$ ), over all possible arrival and error patterns (sets  $A$  and  $E$ ).

$$T_{Alg} = \inf_{A \in \mathcal{A}, E \in \mathcal{E}} \lim_{t \rightarrow \infty} T_{Alg}(A, E, t)$$

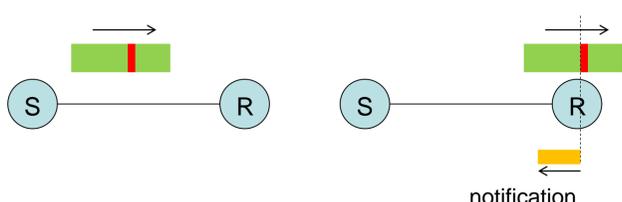
**Observation:** The general offline version of our scheduling problem, where the algorithm has full knowledge of errors and arrivals, is NP-hard.

## Feedback Mechanisms

**Deferred:** The sender is notified of an error after the whole packet has been received by the receiver.



**Instantaneous:** The sender is notified that a packet has suffered an error, at the moment that the error occurs.



This mechanism is an abstraction of the Continuous Error Detection (CED) framework [2].

## One Packet Length

**Theorem 1** No packet scheduling algorithm  $Alg$  can achieve a relative throughput larger than 0 in the deferred feedback model. This holds for both models of packet arrivals; adversarial and stochastic.

**Theorem 2** Any work conserving online scheduling algorithm with instantaneous feedback, has optimal relative throughput of 1 when all packets have the same length.

## Two packet lengths / Instantaneous Feedback

Two different packet lengths,  $l_{min}$  and  $l_{max}$ , where  $l_{min} < l_{max}$ .

Let  $\gamma = l_{max} / l_{min}$  and  $\bar{\gamma} = \lfloor \gamma \rfloor$

### Adversarial Arrivals:

Upper Bound	Lower Bound
$T_{Alg} \leq \bar{\gamma} / (\gamma + \bar{\gamma}) \leq 1/2$	$T_{SL-Pr} \geq \bar{\gamma} / (\gamma + \bar{\gamma})$

### Algorithm SL-Pr:

At the beginning of the execution and whenever the sender is notified about a link error, it checks the queue of pending packets. If there are at least  $\gamma$  packets of length  $l_{min}$ , it schedules them (preamble) and then continues using the *Longest Length* (LL) policy. Otherwise, it simply schedules packets following the LL policy.

### Stochastic Arrivals:

Packets arrive according to a Poisson process with parameter  $\lambda > 0$  and are independently assigned one of the lengths  $l_{min}$  and  $l_{max}$ , with probabilities  $p > 0$  and  $q > 0$  respectively, where  $p + q = 1$ .

Upper Bound	Lower Bound
$T_{Alg} \leq \bar{\gamma} / \gamma$	$T_{CSL-Pr} \geq \bar{\gamma} / (\gamma + \bar{\gamma})$ if $\lambda p l_{min} \leq \bar{\gamma} / (2\gamma)$
$T_{Alg} \leq \max\{\lambda p l_{min}, \bar{\gamma} / (\gamma + \bar{\gamma})\}$ if $p < q$	$T_{CSL-Pr} \geq \min\{\lambda p l_{min}, \bar{\gamma} / \gamma\}$ if $\lambda p l_{min} > \bar{\gamma} / (2\gamma)$

### Algorithm CSL-Pr:

If  $\lambda p l_{min} > \bar{\gamma} / (2\gamma)$ , then algorithm *Shortest Length* (SL) is run, otherwise algorithm SL-Pr is executed.

## References

- [1] Kirk Pruhs, Jiri Sgall, and Eric Torng. *Online scheduling*. pages 115-124. CRC Press, 2003.
- [2] Anand Raghavan, Kannan Ramchandran, and Igor Kozintsev. *Continuous Error Detection (CED) for reliable communication*. IEEE Transactions on Communications, 49(9):1540-1549, 2001.